

A Class of Kazhdan-Lusztig R -Polynomials and q -Fibonacci Numbers

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Abstract

Let S_n denote the symmetric group on $\{1, 2, \dots, n\}$. For two permutations $u, v \in S_n$ such that $u \leq v$ in the Bruhat order, let $R_{u,v}(q)$ and $\tilde{R}_{u,v}(q)$ denote the Kazhdan-Lusztig R -polynomial and \tilde{R} -polynomial, respectively. Let $v_n = 34 \cdots n 12$, and let σ be a permutation such that $\sigma \leq v_n$. We obtain a formula for the R -polynomials $\tilde{R}_{\sigma, v_n}(q)$ in terms of the q -Fibonacci numbers depending on a parameter determined by the reduced expression of σ . When σ is the identity e , this reduces to a formula obtained by Pagliacci. In another direction, we obtain a formula for the \tilde{R} -polynomial $\tilde{R}_{e, v_{n,i}}(q)$, where $v_{n,i} = 34 \cdots i n (i+1) \cdots (n-1) 12$. In a more general context, we conjecture that for any two permutations $\sigma, \tau \in S_n$ such that $\sigma \leq \tau \leq v_n$, the \tilde{R} -polynomial $\tilde{R}_{\sigma, \tau}(q)$ can be expressed as a product of q -Fibonacci numbers multiplied by a power of q .

Keywords: Kazhdan-Lusztig R -polynomial, q -Fibonacci number, symmetric group

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1 Introduction

Let S_n denote the symmetric group on $\{1, 2, \dots, n\}$. For two permutations $u, v \in S_n$ such that $u \leq v$ in the Bruhat order, let $R_{u,v}(q)$ be the Kazhdan-Lusztig R -polynomial, and $\tilde{R}_{u,v}(q)$ be the Kazhdan-Lusztig \tilde{R} -polynomial. Let $v_n = 34 \cdots n 12$, and let σ be a permutation such that $\sigma \leq v_n$. The main result of this paper is a formula for the \tilde{R} -polynomials $\tilde{R}_{\sigma, v_n}(q)$ in terms of the q -Fibonacci numbers depending on a parameter determined by the reduced expression of σ . When σ is the identity permutation e , a formula for the \tilde{R} -polynomials has been given by Pagliacci [6, Theorem 4.1].

We also derive a formula for the \tilde{R} -polynomials $\tilde{R}_{e, v_{n,i}}(q)$, where $v_{n,i} = 34 \cdots i n (i+1) \cdots (n-1) 12$, which can be viewed as a generalization of Pagliacci's formula [6, Theorem 4.1] in another direction. We conclude this paper with a conjecture that for any two permutations $\sigma, \tau \in S_n$ such that $\sigma \leq \tau \leq v_n$, the \tilde{R} -polynomial $\tilde{R}_{\sigma, \tau}(q)$ can be expressed as a product of q -Fibonacci numbers and a power of q .

Let us give an overview of some notation and background. For each permutation π in S_n , it is known that π can be expressed as a product of simple transpositions $s_i = (i, i+1)$ subject

to the following braid relations

$$\begin{aligned} s_i s_j &= s_j s_i, & \text{for } |i - j| > 1; \\ s_i s_{i+1} s_i &= s_{i+1} s_i s_{i+1}, & \text{for } 1 \leq i \leq n - 2. \end{aligned}$$

An expression ω of π is said to be reduced if the number of simple transpositions appearing in ω is minimum. The following word property is due to Tits, see Björner and Brenti [1, Theorem 3.3.1].

Theorem 1.1 (Word Property) *Let π be a permutation of S_n , and ω_1 and ω_2 be two reduced expressions of π . Then ω_1 and ω_2 can be obtained from each other by applying a sequence of braid relations.*

Let $\ell(\pi)$ denote the length of π , that is, the number of simple transpositions in a reduced expression of π . Write $D_R(\pi)$ for the set of right descents of π , namely,

$$D_R(\pi) = \{s_i : 1 \leq i \leq n - 1, \ell(\pi s_i) < \ell(\pi)\}.$$

The exchange condition gives a characterization for the (right) descents of a permutation in terms of reduced expressions, see Humphreys [4, Section 1.7].

Theorem 1.2 (Exchange Condition) *Let $\pi = s_{i_1} s_{i_2} \cdots s_{i_k}$ be a reduced expression of π . If $s_i \in D_R(\pi)$, then there exists an index i_j for which $\pi s_i = s_{i_1} \cdots \widehat{s}_{i_j} \cdots s_{i_k}$, where \widehat{s}_{i_j} means that s_{i_j} is missing. In particular, π has a reduced expression ending with s_i if and only if $s_i \in D_R(\pi)$.*

The following subword property serves as a definition of the Bruhat order. For other equivalent definitions of the Bruhat order, see Björner and Brenti [1]. For a reduced expression $\omega = s_{i_1} s_{i_2} \cdots s_{i_k}$, we say that $s_{i_{j_1}} s_{i_{j_2}} \cdots s_{i_{j_m}}$ is a subword of ω if $1 \leq j_1 < j_2 < \cdots < j_m \leq k$.

Theorem 1.3 (Subword Property) *Let u and v be two permutations in S_n . Then $u \leq v$ in the Bruhat order if and only if every reduced expression of v has a subword that is a reduced expression of u .*

The Bruhat order satisfies the following lifting property, see Björner and Brenti [1, Proposition 2.2.7].

Theorem 1.4 (Lifting Property) *Suppose that u and v are two permutations in S_n such that $u < v$. For any simple transposition s_i in $D_R(v) \setminus D_R(u)$, we have $u \leq v s_i$ and $u s_i \leq v$.*

The Kazhdan-Lusztig R -polynomials, which were introduced by Kazhdan and Lusztig [5], can be recursively determined by the following properties, see also Humphreys [4, Section 7.5].

Theorem 1.5 *For any $u, v \in S_n$,*

- (i) $R_{u,v}(q) = 0$, if $u \not\leq v$;

(ii) $R_{u,v}(q) = 1$, if $u = v$;

(iii) If $u < v$ and $s \in D_R(v)$,

$$R_{u,v}(q) = \begin{cases} R_{us,vs}(q), & \text{if } s \in D_R(u); \\ qR_{us,vs}(q) + (q-1)R_{u,vs}(q), & \text{if } s \notin D_R(u). \end{cases}$$

While R -polynomials may contain negative coefficients, a variant of the R -polynomials introduced by Dyer [3], which has been called the \tilde{R} -polynomials, has only nonnegative coefficients. For an alternative definition of the \tilde{R} -polynomials for the symmetric group, see Brenti [2]. The following two theorems are due to Dyer [3], see also Brenti [2].

Theorem 1.6 *Let $u, v \in S_n$ with $u \leq v$. Then, for $s \in D_R(v)$,*

$$\tilde{R}_{u,v}(q) = \begin{cases} \tilde{R}_{us,vs}(q), & \text{if } s \in D_R(u); \\ \tilde{R}_{us,vs}(q) + q\tilde{R}_{u,vs}(q), & \text{if } s \notin D_R(u). \end{cases} \quad (1.1)$$

Theorem 1.7 *Let $u, v \in S_n$ with $u \leq v$. Then*

$$R_{u,v}(q) = q^{\frac{\ell(v)-\ell(u)}{2}} \tilde{R}_{u,v}(q^{\frac{1}{2}} - q^{-\frac{1}{2}}).$$

Recall that $v_n = 34 \cdots n 12$ and $v_{n,i} = 34 \cdots i n (i+1) \cdots (n-1) 12$. We shall use the recurrence relations in Theorem 1.6 to deduce a formula for the \tilde{R} -polynomials $\tilde{R}_{\sigma, v_n}(q)$, from which we also find a formula for the \tilde{R} -polynomial $\tilde{R}_{e, v_{n,i}}(q)$.

2 Main result

The main result of this paper is an equation for $\tilde{R}_{\sigma, v_n}(q)$, where $v_n = 34 \cdots n 12$ and $\sigma \leq v_n$ in the Bruhat order. Combining this equation with a formula of Pagliacci [6], we obtain an expression of $\tilde{R}_{\sigma, v_n}(q)$ in terms of q -Fibonacci numbers. To describe our result, we need the following reduced expression of v_n .

Proposition 2.1 *For $n \geq 3$,*

$$\Omega_n = s_2 s_1 s_3 s_2 \cdots s_{n-1} s_{n-2}$$

is a reduced expression of v_n .

Let σ be a permutation of S_n such that $\sigma \leq v_n$. By the subword property in Theorem 1.3, σ can be expressed as a reduced subword of Ω_n . We introduce two statistics of a reduced subword of Ω_n .

Let $\omega = s_{i_1} s_{i_2} \cdots s_{i_k}$ be a reduced subword of Ω_n . Define

$$D(\omega) = \{1 \leq t < k : i_t - i_{t+1} = 1\}.$$

We use $d(\omega)$ to denote the cardinality of $D(\omega)$, and let

$$h(\omega) = n - \ell(\omega) + d(\omega). \quad (2.1)$$

For example, for a reduced subword $\omega = s_2s_3s_4s_3s_6s_5s_7$ of Ω_9 , we have $D(\omega) = \{3, 5\}$, and thus $d(\omega) = 2$ and $h(\omega) = 4$. Note that $h(\omega)$ depends on both ω and n . This causes no confusion since the index n is always clear from the context.

The main result in this paper is the following equation for the \tilde{R} -polynomials $\tilde{R}_{\sigma, v_n}(q)$.

Theorem 2.2 *For $n \geq 3$, let σ be a permutation in S_n such that $\sigma \leq v_n$, and let ω be any reduced expression of σ that is a subword of Ω_n . Then we have*

$$\tilde{R}_{\sigma, v_n}(q) = q^{\ell(\omega) - 2d(\omega)} \tilde{R}_{e, v_{h(\omega)}}(q). \quad (2.2)$$

Let $F_n(q)$ be the q -Fibonacci numbers, that is, $F_0(q) = F_1(q) = 1$ and for $n \geq 2$,

$$F_n(q) = F_{n-1}(q) + qF_{n-2}(q).$$

Pagliacci [6, Theorem 4.1] has shown that

$$\tilde{R}_{e, v_n}(q) = q^{2n-4} F_{n-2}(q^{-2}). \quad (2.3)$$

As a consequence of Theorem 2.2 and formula (2.3), we obtain an expression of $\tilde{R}_{\sigma, v_n}(q)$ in terms of q -Fibonacci numbers.

Corollary 2.3 *For $n \geq 3$, let σ be a permutation in S_n such that $\sigma \leq v_n$, and let ω be any reduced expression of σ that is a subword of Ω_n . Then we have*

$$\tilde{R}_{\sigma, v_n}(q) = q^{2n - \ell(\sigma) - 4} F_{h(\omega) - 2}(q^{-2}). \quad (2.4)$$

To give an inductive proof of Theorem 2.2, we need three lemmas. Assume that ω is a reduced subword of Ω_n . When $\omega s_{n-1} \leq \Omega_n$, the first two lemmas are concerned with the existence of a reduced expression ω' of ωs_{n-1} such that $d(\omega') = d(\omega)$. When $\omega s_{n-1} \not\leq \Omega_n$, the third lemma shows that $h(\omega) = 2$.

Lemma 2.4 *Let ω be a reduced subword of Ω_n . If $\omega s_{n-1} \leq \Omega_n$ and $s_{n-1} \in D_R(\omega)$, then there exists a reduced expression ω' of ωs_{n-1} such that ω' is a subword of Ω_n and $d(\omega') = d(\omega)$.*

Proof. We use induction on n . It is easy to check that the lemma holds for $n \leq 3$. Assume that $n > 3$ and the assertion holds for $n - 1$. We now consider the case for n . By definition, we have $\Omega_n = \Omega_{n-1} s_{n-1} s_{n-2}$. Since ω is a subword of Ω_n , we can write $\omega = \omega_1 \omega_2$, where ω_1 is a subword of Ω_{n-1} and ω_2 is a subword of $s_{n-1} s_{n-2}$. Because $s_{n-1} \in D_R(\omega)$, we have the following two cases.

Case 1: $\omega = \omega_1 s_{n-1}$. Set $\omega' = \omega_1$. Clearly, ω' is a reduced expression of ωs_{n-1} . Moreover, it is easy to check that $D(\omega) = D(\omega')$, and thus $d(\omega') = d(\omega)$, that is, ω' is a desired reduced expression of ωs_{n-1} .

Case 2: $\omega = \omega_1 s_{n-1} s_{n-2}$. Since $s_{n-1} \in D_R(\omega)$, by Theorem 1.2, there exists a reduced expression of ω ending with s_{n-1} . Hence the word property in Theorem 1.1 ensures the existence

of a reduced expression of ω_1 ending with s_{n-2} . This implies that s_{n-2} belongs to $D_R(\omega_1)$. By the induction hypothesis, there exists a reduced expression ω'_1 of $\omega_1 s_{n-2}$ such that ω'_1 is a subword of Ω_{n-1} and $d(\omega'_1) = d(\omega_1)$.

Set

$$\omega' = \omega'_1 s_{n-1} s_{n-2}.$$

We deduce that ω' is a desired reduced subword. Since

$$\omega' = \omega'_1 s_{n-1} s_{n-2} = \omega_1 s_{n-2} s_{n-1} s_{n-2} = \omega_1 s_{n-1} s_{n-2} s_{n-1} = \omega s_{n-1},$$

we see that ω' is an expression of ωs_{n-1} . On the other hand, since ω' consists of $\ell(\omega'_1) + 2$ simple transpositions and

$$\ell(\omega'_1) + 2 = \ell(\omega_1) + 1 = \ell(\omega) - 1 = \ell(\omega s_{n-1}),$$

we conclude that ω' is a reduced expression of ωs_{n-1} . By the construction of ω' , we have

$$d(\omega') = d(\omega'_1) + 1 = d(\omega_1) + 1 = d(\omega).$$

This completes the proof. ■

The next lemma deals with the case $s_{n-1} \notin D_R(\omega)$.

Lemma 2.5 *Let ω be a reduced subword of Ω_n . If $\omega s_{n-1} \leq \Omega_n$ and $s_{n-1} \notin D_R(\omega)$, then there exists a reduced expression ω' of ωs_{n-1} such that ω' is a subword of Ω_n and $d(\omega') = d(\omega)$.*

Proof. We use induction on n . It is easily checked that the lemma holds for $n \leq 3$. Assume that $n > 3$ and the assertion holds for $n - 1$. We now consider the case for n . Let $\omega = \omega_1 \omega_2$, where ω_1 is a subword of Ω_{n-1} and ω_2 is a subword of $s_{n-1} s_{n-2}$. Since $s_{n-1} \notin D_R(\omega)$, we have the following three cases.

Case 1: $\omega = \omega_1$. Set $\omega' = \omega_1 s_{n-1}$. It is easily seen that ω' is a desired reduced expression.

Case 2: $\omega = \omega_1 s_{n-2}$. We claim that $\omega = \omega_1 s_{n-2} \leq \Omega_{n-1}$. Note that s_{n-1} does not appear in ω_1 . Since $\omega s_{n-1} = \omega_1 s_{n-2} s_{n-1}$, by Theorem 1.1, there does not exist any reduced expression of ωs_{n-1} ending with s_{n-2} . This implies that s_{n-2} does not belong to $D_R(\omega s_{n-1})$. Thus, by the lifting property in Theorem 1.4, we deduce that

$$\omega s_{n-1} \leq \Omega_n s_{n-2} = \Omega_{n-1} s_{n-1}.$$

This implies that $\omega = \omega_1 s_{n-2} \leq \Omega_{n-1}$, as claimed.

Since $\omega = \omega_1 s_{n-2}$ is reduced, we see that $s_{n-2} \notin D_R(\omega_1)$. By the induction hypothesis, there exists a reduced expression ω'_1 of $\omega_1 s_{n-2}$ such that ω'_1 is a subword of Ω_{n-1} and $d(\omega'_1) = d(\omega_1)$. Set $\omega' = \omega'_1 s_{n-1}$. We find that ω' is a reduced expression of ωs_{n-1} such that $d(\omega') = d(\omega)$.

Case 3: $\omega = \omega_1 s_{n-1} s_{n-2}$. We claim that $s_{n-2} \notin D_R(\omega_1)$. Suppose to the contrary that $s_{n-2} \in D_R(\omega_1)$. By Theorem 1.2, there exists a reduced expression of ω_1 ending with s_{n-2} . Write $\omega_1 = \mu s_{n-2}$, where μ is a reduced expression. Then we get

$$\omega = \mu s_{n-2} s_{n-1} s_{n-2} = \mu s_{n-1} s_{n-2} s_{n-1},$$

contradicting the assumption that $s_{n-1} \notin D_R(\omega)$. So the claim is proved.

On the other hand, since

$$\omega_1 s_{n-2} s_{n-1} s_{n-2} = \omega_1 s_{n-1} s_{n-2} s_{n-1} = \omega s_{n-1} \leq \Omega_n,$$

we have $\omega_1 s_{n-2} \leq \Omega_{n-1}$. It follows from the induction hypothesis that there exists a reduced expression ω'_1 of $\omega_1 s_{n-2}$ such that ω'_1 is a subword of Ω_{n-1} and $d(\omega'_1) = d(\omega_1)$.

Let

$$\omega' = \omega'_1 s_{n-1} s_{n-2}.$$

Since

$$\omega' = \omega'_1 s_{n-1} s_{n-2} = \omega_1 s_{n-2} s_{n-1} s_{n-2} = \omega_1 s_{n-1} s_{n-2} s_{n-1} = \omega s_{n-1},$$

we deduce that ω' is a reduced expression of ωs_{n-1} . By the construction of ω' , we obtain that

$$d(\omega') = d(\omega'_1) + 1 = d(\omega_1) + 1 = d(\omega),$$

as required. ■

We now come to the third lemma.

Lemma 2.6 *Let ω be a reduced subword of Ω_n . If $\omega s_{n-1} \not\leq \Omega_n$, then we have $h(\omega) = 2$.*

Proof. We proceed by induction on n . It can be verified that the lemma holds for $n \leq 3$. Assume that $n > 3$ and the assertion holds for $n - 1$. Consider the case for n . Write $\omega = \omega_1 \omega_2$, where ω_1 is a subword of Ω_{n-1} and ω_2 is a subword of $s_{n-1} s_{n-2}$. Since $\omega s_{n-1} \not\leq \Omega_n$, we see that s_{n-1} is not a right descent of ω . We have the following two cases.

Case 1: $\omega = \omega_1 s_{n-2}$. Since $\omega s_{n-1} \not\leq \Omega_n$, we have $\omega = \omega_1 s_{n-2} \not\leq \Omega_{n-1}$. Thus, by the induction hypothesis, we get $h(\omega_1) = 2$. Noticing that $\ell(\omega_1) = \ell(\omega) - 1$ and $d(\omega_1) = d(\omega)$, we obtain that

$$h(\omega) = n - \ell(\omega) + d(\omega) = n - 1 - \ell(\omega_1) + d(\omega_1) = h(\omega_1) = 2,$$

as required.

Case 2: $\omega = \omega_1 s_{n-1} s_{n-2}$. We claim that $\omega_1 s_{n-2} \not\leq \Omega_{n-1}$. Suppose to the contrary that $\omega_1 s_{n-2} \leq \Omega_{n-1}$. Note that

$$\omega s_{n-1} = \omega_1 s_{n-1} s_{n-2} s_{n-1} = \omega_1 s_{n-2} s_{n-1} s_{n-2}.$$

This yields $\omega s_{n-1} \leq \Omega_n$, contradicting the assumption that $\omega s_{n-1} \not\leq \Omega_n$. So the claim is verified.

By the induction hypothesis, we have $h(\omega_1) = 2$. Since $\ell(\omega_1) = \ell(\omega) - 2$ and $d(\omega_1) = d(\omega) - 1$, we find that

$$h(\omega) = n - \ell(\omega) + d(\omega) = n - 1 - \ell(\omega_1) + d(\omega_1) = h(\omega_1) = 2,$$

as required. ■

Next we give a proof of Theorem 2.2 based on the above lemmas.

Proof of Theorem 2.2. Let

$$T_n(q) = \tilde{R}_{e, v_n}(q).$$

and

$$g(\omega) = \ell(\omega) - 2d(\omega). \tag{2.5}$$

Then equation (2.2) can be rewritten as

$$\tilde{R}_{\sigma, v_n}(q) = q^{g(\omega)} T_{h(\omega)}(q). \quad (2.6)$$

We proceed to prove (2.6) by induction on n . It can be checked that (2.6) holds for $n \leq 3$. Assume that $n > 3$ and (2.6) holds for $n - 1$. For the case for n , let $\omega = \omega_1 \omega_2$, where ω_1 is a subword of Ω_{n-1} and ω_2 is a subword of $s_{n-1} s_{n-2}$. There are four cases.

Case 1: $\omega = \omega_1 s_{n-2}$. It follows from (1.1) that

$$\begin{aligned} \tilde{R}_{\sigma, v_n}(q) &= \tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1} s_{n-1} s_{n-2}}(q) \\ &= \tilde{R}_{\omega_1, \Omega_{n-1} s_{n-1}}(q) \\ &= \tilde{R}_{\omega_1 s_{n-1}, \Omega_{n-1}}(q) + q \tilde{R}_{\omega_1, \Omega_{n-1}}(q). \end{aligned} \quad (2.7)$$

Since $\omega_1 s_{n-1} \not\leq \Omega_{n-1}$, we see that the first term in (2.7) vanishes. Thus, (2.7) becomes

$$\tilde{R}_{\sigma, v_n}(q) = q \tilde{R}_{\omega_1, \Omega_{n-1}}(q). \quad (2.8)$$

By the induction hypothesis, we have

$$\tilde{R}_{\omega_1, \Omega_{n-1}}(q) = q^{g(\omega_1)} T_{h(\omega_1)}. \quad (2.9)$$

Since $d(\omega_1) = d(\omega)$, we get

$$g(\omega_1) = \ell(\omega_1) - 2d(\omega_1) = \ell(\omega) - 1 - 2d(\omega) = g(\omega) - 1 \quad (2.10)$$

and

$$h(\omega_1) = n - 1 - \ell(\omega_1) + d(\omega_1) = n - \ell(\omega) + d(\omega) = h(\omega). \quad (2.11)$$

Plugging (2.10) and (2.11) into (2.9), we obtain

$$\tilde{R}_{\omega_1, \Omega_{n-1}}(q) = q^{g(\omega_1)} T_{h(\omega_1)} = q^{g(\omega)-1} T_{h(\omega)}(q),$$

which leads to

$$\tilde{R}_{\sigma, v_n}(q) = q \tilde{R}_{\omega_1, \Omega_{n-1}}(q) = q^{g(\omega)} T_{h(\omega)}(q).$$

Case 2: $\omega = \omega_1 s_{n-1}$. By Theorem 1.1, there is no reduced expression of σ that ends with s_{n-2} . This implies that s_{n-2} is not a right descent of σ , so that

$$\begin{aligned} \tilde{R}_{\sigma, v_n}(q) &= \tilde{R}_{\omega_1 s_{n-1} s_{n-2}, \Omega_{n-1} s_{n-1}}(q) + q \tilde{R}_{\omega_1 s_{n-1}, \Omega_{n-1} s_{n-1}}(q) \\ &= \tilde{R}_{\omega_1 s_{n-1} s_{n-2}, \Omega_{n-1} s_{n-1}}(q) + q \tilde{R}_{\omega_1, \Omega_{n-1}}(q). \end{aligned} \quad (2.12)$$

We claim that the first term in (2.12) vanishes, or equivalently, $\omega_1 s_{n-1} s_{n-2} \not\leq \Omega_{n-1} s_{n-1}$. Suppose to the contrary that $\omega_1 s_{n-1} s_{n-2} \leq \Omega_{n-1} s_{n-1}$. By Theorem 1.3, there exists a subword μ of $\Omega_{n-1} s_{n-1}$ that is a reduced expression of $\omega_1 s_{n-1} s_{n-2}$. Since s_{n-1} must appear in μ , we may write μ in the following form

$$\mu = s_{i_1} s_{i_2} \cdots s_{i_k} s_{n-1},$$

where $s_{i_1} s_{i_2} \cdots s_{i_k}$ is a reduced subword of Ω_{n-1} . By the word property in Theorem 1.1, $\omega_1 s_{n-1} s_{n-2}$ can be obtained from μ by applying the braid relations. However, this is impossible

since any simple transposition s_{n-2} appearing in μ cannot be moved to the last position by applying the braid relations. So the claim is proved, and hence (2.12) becomes

$$\tilde{R}_{\sigma, v_n}(q) = q\tilde{R}_{\omega_1, \Omega_{n-1}}(q).$$

It is easily seen that

$$g(\omega_1) = g(\omega) - 1 \quad \text{and} \quad h(\omega_1) = h(\omega).$$

By the induction hypothesis, we deduce that

$$\tilde{R}_{\sigma, v_n}(q) = q\tilde{R}_{\omega_1, \Omega_{n-1}}(q) = q^{g(\omega_1)+1}T_{h(\omega_1)}(q) = q^{g(\omega)}T_{h(\omega)}(q).$$

Case 3: $\omega = \omega_1 s_{n-1} s_{n-2}$. It is clear from (1.1) that

$$\tilde{R}_{\sigma, v_n}(q) = \tilde{R}_{\omega_1 s_{n-1}, \Omega_{n-1} s_{n-1}}(q) = \tilde{R}_{\omega_1, \Omega_{n-1}}(q). \quad (2.13)$$

Noting that $d(\omega) = d(\omega_1) + 1$, we obtain

$$g(\omega_1) = \ell(\omega_1) - 2d(\omega_1) = \ell(\omega_1) + 2 - 2d(\omega) = \ell(\omega) - 2d(\omega) = g(\omega)$$

and

$$h(\omega_1) = n - 1 - \ell(\omega_1) + d(\omega_1) = n - 2 - \ell(\omega_1) + d(\omega) = n - \ell(\omega) + d(\omega) = h(\omega).$$

Thus, by (2.13) and the induction hypothesis, we find that

$$\tilde{R}_{\sigma, v_n}(q) = \tilde{R}_{\omega_1, \Omega_{n-1}}(q) = q^{g(\omega_1)}T_{h(\omega_1)}(q) = q^{g(\omega)}T_{h(\omega)}(q).$$

Case 4: $\omega = \omega_1$. Here are two subcases.

Subcase 1: $s_{n-2} \in D_R(\omega_1)$. By (1.1), we deduce that

$$\begin{aligned} \tilde{R}_{\sigma, v_n}(q) &= \tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1} s_{n-1}}(q) \\ &= q\tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1}}(q). \end{aligned} \quad (2.14)$$

By Lemma 2.4, there exists a reduced expression ω'_1 of $\omega_1 s_{n-2}$ such that ω'_1 is a subword of Ω_{n-1} and $d(\omega'_1) = d(\omega_1)$. Consequently,

$$g(\omega'_1) = \ell(\omega'_1) - 2d(\omega'_1) = \ell(\omega_1) - 1 - 2d(\omega_1) = g(\omega_1) - 1$$

and

$$h(\omega'_1) = n - 1 - \ell(\omega'_1) + d(\omega'_1) = n - \ell(\omega_1) + d(\omega_1) = h(\omega_1) + 1.$$

By the induction hypothesis, we obtain that

$$\begin{aligned} \tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1}}(q) &= \tilde{R}_{\omega'_1, \Omega_{n-1}}(q) \\ &= q^{g(\omega'_1)}T_{h(\omega'_1)}(q) \\ &= q^{g(\omega_1)-1}T_{h(\omega_1)+1}(q). \end{aligned} \quad (2.15)$$

But $h(\omega) = h(\omega_1) + 1$, substituting (2.15) into (2.14) gives

$$\begin{aligned} \tilde{R}_{\sigma, v_n}(q) &= q\tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1}}(q) \\ &= q^{g(\omega)}T_{h(\omega)}(q). \end{aligned}$$

Subcase 2: $s_{n-2} \notin D_R(\omega_1)$. By (1.1), we see that

$$\begin{aligned}\tilde{R}_{\sigma, v_n}(q) &= \tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1} s_{n-1}}(q) + q \tilde{R}_{\omega_1, \Omega_{n-1} s_{n-1}}(q) \\ &= q \tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1}}(q) + q^2 \tilde{R}_{\omega_1, \Omega_{n-1}}(q)\end{aligned}\quad (2.16)$$

By the induction hypothesis, the second term in (2.16) equals

$$q^2 \tilde{R}_{\omega_1, \Omega_{n-1}}(q) = q^{g(\omega_1)+2} T_{h(\omega_1)}(q) = q^{g(\omega)+2} T_{h(\omega)-1}(q). \quad (2.17)$$

It remains to compute the first term in (2.16). To this end, we have the following two cases.

Subcase 2a: $\omega_1 s_{n-2} \leq \Omega_{n-1}$. By Lemma 2.5, there exists a reduced expression ω'_1 of $\omega_1 s_{n-2}$ such that ω'_1 is a subword of Ω_{n-1} and $d(\omega_1) = d(\omega'_1)$. Hence

$$g(\omega'_1) = \ell(\omega'_1) - 2d(\omega'_1) = \ell(\omega_1) + 1 - 2d(\omega_1) = g(\omega_1) + 1 = g(\omega) + 1$$

and

$$h(\omega'_1) = n - 1 - \ell(\omega'_1) + d(\omega'_1) = n - 2 - \ell(\omega_1) + d(\omega_1) = h(\omega_1) - 1 = h(\omega) - 2.$$

By the induction hypothesis, we obtain that

$$\tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1}}(q) = q^{g(\omega'_1)} T_{h(\omega'_1)}(q) = q^{g(\omega)+1} T_{h(\omega)-2}(q). \quad (2.18)$$

Putting (2.17) and (2.18) into (2.16), we deduce that

$$\begin{aligned}\tilde{R}_{\sigma, v_n}(q) &= q \tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1}}(q) + q^2 \tilde{R}_{\omega_1, \Omega_{n-1}}(q) \\ &= q^{g(\omega)+2} T_{h(\omega)-2}(q) + q^{g(\omega)+2} T_{h(\omega)-1}(q) \\ &= q^{g(\omega)} (q^2 T_{h(\omega)-2}(q) + q^2 T_{h(\omega)-1}(q)).\end{aligned}\quad (2.19)$$

In view of the following relation due to Pagliacci [6]

$$T_n(q) = q^2 T_{n-2}(q) + q^2 T_{n-1}(q),$$

(2.19) can be rewritten as

$$\tilde{R}_{\sigma, v_n}(q) = q^{g(\omega)} T_{h(\omega)}(q).$$

Subcase 2b: $\omega_1 s_{n-2} \not\leq \Omega_{n-1}$. In this case, we have

$$\tilde{R}_{\omega_1 s_{n-2}, \Omega_{n-1}}(q) = 0. \quad (2.20)$$

By Lemma 2.6, we find that $h(\omega_1) = 2$. Thus (2.17) reduces to

$$q^2 \tilde{R}_{\omega_1, \Omega_{n-1}}(q) = q^{g(\omega)+2} T_2(q) = q^{g(\omega)+2}. \quad (2.21)$$

Putting (2.20) and (2.21) into (2.16), we get

$$\tilde{R}_{\sigma, v_n}(q) = q^{g(\omega)+2}. \quad (2.22)$$

Since $T_3(q) = q^2$ and $h(\omega) = h(\omega_1) + 1 = 3$, it follows from (2.22) that

$$\tilde{R}_{\sigma, v_n}(q) = q^{g(\omega)} T_3(q) = q^{g(\omega)} T_{h(\omega)}(q),$$

and hence the proof is complete. ■

For $2 \leq i \leq n-1$, let

$$v_{n,i} = \begin{cases} n34 \cdots (n-1)12, & \text{if } i = 2; \\ 34 \cdots i n (i+1) \cdots (n-1)12, & \text{if } 3 \leq i \leq n-1. \end{cases}$$

We obtain the following formula for $\tilde{R}_{e, v_{n,i}}(q)$, which reduces to formula (2.3) due to Pagliacci in the case $i = n-1$.

Theorem 2.7 *Let $n \geq 3$ and $2 \leq i \leq n-1$. Then we have*

$$\tilde{R}_{e, v_{n,i}}(q) = \sum_{k=0}^{n-i-1} q^{3n-i-2k-5} \binom{n-i-1}{k} F_{n-k-2}(q^{-2}). \quad (2.23)$$

Proof. Recall that

$$T_n(q) = \tilde{R}_{e, v_n}(q) = q^{2n-4} F_{n-2}(q^{-2}),$$

see (2.3). Hence (2.23) can be rewritten as

$$\tilde{R}_{e, v_{n,i}}(q) = \sum_{k=0}^{n-i-1} q^{n-i-1} \binom{n-i-1}{k} T_{n-k}(q). \quad (2.24)$$

Note that $\Omega_n s_{n-3} s_{n-4} \cdots s_{i-1}$ is a reduced expression of the permutation $v_{n,i}$. By the defining relation (1.1) of \tilde{R} -polynomials, we obtain that

$$\begin{aligned} \tilde{R}_{e, v_{n,i}}(q) &= \tilde{R}_{e, \Omega_n s_{n-3} s_{n-4} \cdots s_{i-1}}(q) \\ &= \tilde{R}_{s_{i-1}, \Omega_n s_{n-3} s_{n-4} \cdots s_i}(q) + q \tilde{R}_{e, \Omega_n s_{n-3} s_{n-4} \cdots s_i}(q) \\ &= \left(\tilde{R}_{s_{i-1} s_i, \Omega_n s_{n-3} s_{n-4} \cdots s_{i+1}}(q) + q \tilde{R}_{s_{i-1}, \Omega_n s_{n-3} s_{n-4} \cdots s_{i+1}}(q) \right) \\ &\quad + q \left(\tilde{R}_{s_i, \Omega_n s_{n-3} s_{n-4} \cdots s_{i+1}}(q) + q \tilde{R}_{e, \Omega_n s_{n-3} s_{n-4} \cdots s_{i+1}}(q) \right) \\ &= \cdots \\ &= \sum_{i-1 \leq i_1 < \cdots < i_k \leq n-3} q^{n-i-1-k} \tilde{R}_{s_{i_1} \cdots s_{i_k}, \Omega_n}(q). \end{aligned} \quad (2.25)$$

Observe that $s_{i_1} \cdots s_{i_k}$ is a reduced subword of Ω_n with $d(s_{i_1} \cdots s_{i_k}) = 0$. By Corollary 2.3, we find that

$$\tilde{R}_{s_{i_1} \cdots s_{i_k}, \Omega_n}(q) = q^k T_{n-k}(q). \quad (2.26)$$

Substituting (2.26) into (2.25), we get

$$\begin{aligned} \tilde{R}_{e, v_{n,i}}(q) &= \sum_{i-1 \leq i_1 < \cdots < i_k \leq n-3} q^{n-i-1-k} \tilde{R}_{s_{i_1} \cdots s_{i_k}, \Omega_n}(q) \\ &= \sum_{i-1 \leq i_1 < \cdots < i_k \leq n-3} q^{n-i-1} T_{n-k}(q) \\ &= \sum_{k=0}^{n-i-1} q^{n-i-1} \binom{n-i-1}{k} T_{n-k}(q), \end{aligned}$$

as required. ■

We conclude this paper with the following conjecture, which has been verified for $n \leq 9$.

Conjecture 2.8 For $n \geq 2$ and $e \leq \sigma_1 \leq \sigma_2 \leq \Omega_n$, we have

$$\tilde{R}_{\sigma_1, \sigma_2}(q) = q^{g(\sigma_1, \sigma_2)} \prod_{i=1}^k F_{h_i(\sigma_1, \sigma_2)}(q^{-2}), \quad (2.27)$$

where k , $g(\sigma_1, \sigma_2)$ and $h_i(\sigma_1, \sigma_2)$ are integers depending on σ_1 and σ_2 .

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